Translating Haskell to Isabelle

Paolo Torrini[1] and Christoph Lueth[2] Christian Maeder[2] Till Mossakowski[2]

[1] Verimag, Grenoble, `Paolo.Torrini@imag.fr`
[2] DFKI Lab Bremen,
{`Christoph.Lueth,Christian.Maeder,Till.Mossakowski`}`@dfki.de`

Abstract. We present partial translations of Haskell programs to Isabelle that have been implemented as part of the Heterogenous Tool Set. The the target logic is Isabelle/HOLCF, and the translation is based on a shallow embedding approach.

1 Introduction

Automating the translation from programming languages to the languagee of a generic prover may provide useful support for the formal development and the verification of programs. It has been argued that functional languages can make the task of proving assertions about programs written in them easier, owing to the relative simplicity of their semantics [Tho92,Tho94]. The idea of translating Haskell programs, came to us, more specifically, from an interest in the use of functional languages for the specification of reactive systems. Haskell is a strongly typed, purely functional language with lazy evaluation, polymorphic types extended with type constructor classes, and a syntax for side effects and pseudo-imperative code based on monadic operators [PJ03]. Several languages based on Haskell have been proposed for application to robotics [PHH99,HCNP03]. In such languages, monadic constructors are extensively used to deal with side-effects. Isabelle is a generic theorem-prover implemented in SML supporting several logics — in particular, Isabelle/HOL is the implementation in Isabelle of classical higher-order logic based on simply typed lambda calculus extended with axiomatic type classes. It provides support for reasoning about programming functions, both in terms of rich libraries and efficient automation. Isabelle/HOLCF [MNvOS99] [Pau94,MNvOS99] is Isabelle/HOL conservatively extended with the Logic of Computable Functions — a formalisation of domain theory.

We have implemented as functions of Hets translations of Haskell to Isabelle/HOLCF following an approach based on shallow embedding, mapping Haskell types to Isabelle ones, therefore taking full advantage of Isabelle builtin type-checking. Hets [Mos05a,Mos06,MML07] is an Haskell-based application designed to support heterogeneous specification and the formal development of programs. It has an interface with Isabelle, and relies on Programatica [HHJK04] for parsing and static analysis of Haskell programs. Programatica already includes a translation to Isabelle/HOLCF which, in contrast to ours, is based on an object-level modelling of the type system [HMW05].

The paper is organised as follows: Section 2 discusses the semantic background of the translation, while the subsequent sections are devoted to the translation of types, datatypes, classes and function definitions, respectively.

2 Semantic Background of the Translation

We firstly describe the subset of Haskell that we cover. Our translation to Isabelle/HOLCF covers at present Booleans, integers, basic constructors (function, product, list, *maybe*), equality, single-parameter type classes (with some limitations), *case* and *if* expressions, *let* expressions without patterns, mutually recursive data-types and functions. It keeps into account partiality and laziness by following, for the main lines, the denotational semantics of lazy evaluation given in [Win93]. There are several limitations: *Predulde* syntax is covered only partially; list comprension, *where* expressions and *let* with patterns are not covered; further built-in types and type classes are not covered; imports are not allowed; constructor type classes are not covered at all — and so for monadic types beyond list and *maybe*. Of all these limitations, the only logically deep ones are those related to classes — all the other ones are just a matter of implementation.

For the translation, we have followed the informal description of the operational semantics given in the Haskell report [PJ03]. However, it should be noted that there is no denotational semantics of Haskell! This also has been confirmed in the answers to a query that one of the authors has sent to the Haskell mailing list. Hence, our translation to Isabelle/HOLCF can be seen as the first denotational semantics given to a large subset of Haskell 98. This also means that there is no notion of correctness of this translation, because it just *defines* the denotational semantics. Of course, an interesting question is the coincidence of denotational and operational semantics. However, this is beyond the scope of the paper.

3 Translation of Types

In Isabelle/HOL types are interpreted as sets (class *type*); functions are total and may not be computable. A non-primitive recursive function may require discharging proof obligations already at the stage of definition — in fact, a specific relation has to be given for a function to be proved total. In Isabelle/HOLCF each type is interpreted as a pointed complete partially ordered set (class *pcpo*) i.e. a set with a partial order which has suprema of ω-chains and has a bottom. Isabelle's formalisation, based on axiomatic type classes [Wen05], makes it possible to deal with complete partial orders in quite an abstract way. Functions are generally partial and computability can be expressed in terms of continuity. Recursion can be expressed in terms of least fixed-point operator, and so, in contrast with Isabelle/HOL, function definition does not depend on proofs. Nevertheless, proving theorems in Isabelle/HOLCF may turn out to be comparatively hard. After being spared the need to discharge proof obligations at

the definition stage, one has to bear with assumptions on function continuity throughout the proofs. A standard strategy is then to define as much as possible in Isabelle/HOL, using Isabelle/HOLCF type constructors to lift types only when this is necessary.

The provision of pcpos, domains and continuous functions by Isabelle/HOLCF eases the translation of Haskell types and functions a lot. However, special care is needed when trying to understand the Haskell semantics. If one reads the section of the Haskell report dealing with types and classes, one could come to the conclusion that a function space in Haskell can be mapped to the space of the continuous functions in Isabelle/HOLCF — this would correspond to a purely lazy language. However, Haskell is a mixed eager and lazy language, and it provides a function *seq* that enforces eager evaluation. This function is introduced in part 6 of the Haskell report, "Predefined Types and Classes", in section 6.2. We quote from there:

> However, the provision of *seq* has important semantic consequences, because it is available at every type. As a consequence, $\perp$ is not the same as $\lambda x \rightarrow \perp$, since *seq* can be used to distinguish them.

In order to enforce this distinction, each function space needs to be lifted. The same holds for products. We define these liftings in a specific Isabelle theory *HsHOLCF* (included in the Hets distrubution) as follows:

```
defaultsort pcpo

domain ('a,'b) lprod = lpair (lazy lfst :: 'a) (lazy lsnd :: 'b)

domain 'a Lift = Lift (lazy 'a)

types  ('a, 'b) "-->"    = "('a -> 'b) Lift" (infixr 0)

constdefs
  liftedApp :: "('a --> 'b) => ('a => 'b)" ("_$$_" [999,1000] 999)
                                              (* application *)
  "liftedApp f x == case f of
                      Lift $ g => g $ x"
constdefs
  liftedLam :: "('a => 'b) => ('a --> 'b)" (binder "Lam " 10)
                                              (* abstraction *)
  "liftedLam f == Lift $ (LAM x . f x)"
```

Our translation of Haskell types to Isabelle types is defined recursively. It is based on a translation of names for avoidance of name clashes that is not specified here. We write α' for both the recursive translation of item α and the renaming according to the name translation. The translation of types is given by the following rules:

Types

$$
\begin{array}{ll}
a & \implies\ 'a :: \{pcpo\} \\
() & \implies\ unit\ \text{lift} \\
Bool & \implies\ bool\ \text{lift} \\
Integer & \implies\ int\ \text{lift} \\
\tau_1 \rightarrow \tau_2 & \implies\ \tau_1' \ \text{-->}\ \tau_2') \\
(\tau_1, \tau_2) & \implies\ (\tau_1'\ lprod\ \tau_2') \\
[\tau] & \implies\ \tau'\ llist \\
T\ \tau_1\ \ldots\ \tau_n & \implies\ \tau_1'\ \ldots\ \tau_n'\ T'
\end{array}
$$

with T either datatype or defined type

Built-in types are translated to the lifting of ths corresponding HOL type. The Isabelle/HOLCF type constructor *lift* is used to lift types to flat domains. The type constructor *llist* is dicussed in the next section.

4 Translation of Datatypes

As explained in the Haskell report [PJ03], section 4.2.3, the following four Haskell declarations

```
data D1 = D1 Int
data D2 = D2 !Int
type S = Int
newtype N = N Int
```

have four different semantics. Indeed, the correct translation to Isabelle/HOLCF is as follows:

```
domain D1 = D1 (lazy D1_1::"Int \mathit{lift}")
domain D2 = D2 (D2_1::"Int \mathit{lift}")
types S = "Int \mathit{lift}"
pcpodef N = "{x:: Int \mathit{lift} . True}"
by auto
```

In Isabelle/HOLCF, the keyword *domain* defines a (possibly recursive) domain as solution of the corresponding domain equation. The keyword *lazy* ensures that the constructor $D1$ is non-strict, i.e. $D1 \perp\ \neq\ \perp$. The keyword *pcpodef* can be used to define sub-pcpos of existing pcpos; here, we use it just to introduce an isomorphic copy of an existing pcpo — this is the semantic of Haskell *newtype* definitions.

Lists are translated to the domain *llist*, defined as follows in our prelude theory *HsHOLCF*:

$domain\ 'a\ llist\ =\ nil\ |\ \#\#\ (lazy\ HD :: \ 'a)\ (lazy\ TL :: \ 'a\ llist)$

allowing for partial list as well as for infinite ones [MNvOS99].

The general scheme for translation of mutually recursive lazy Haskell datatypes to Isabelle/HOLCF domains is as follows:

$$data \ \phi_n \ = \ C_{n1} \ w_1 \ldots w_h \ | \ \ldots \ | \ C_{1q} \ z_1 \ldots z_k \implies$$
$$domain \ \phi'_1 \ = \ C'_{11} \ (lazy \ d_{111} :: \ x'_1) \ \ldots \ (lazy \ d_{11i} :: \ x'_i) \ |$$
$$\ldots \ |$$
$$C'_{1p} \ (lazy \ d_{1p1} :: \ y'_1) \ \ldots \ (lazy \ d_{1pj} :: \ y'_j)$$
$$and \ \ldots$$
$$and \quad \phi'_n \ = \ C'_{n1} \ (lazy \ d_{n11} :: \ w'_1) \ \ldots \ (lazy \ d_{n1h} \ :: \ w'_h) \ |$$
$$\ldots \ |$$
$$C'_{nq} \ (lazy \ d_{nq1} :: \ z'_1) \ \ldots \ (lazy \ d_{nqk} :: \ z'_k)$$

Here, the d_{nqk} are the selectors.

5   Translation of Kinds and Type Classes

Classes
$$Eq \implies Eq$$
$$K \ \implies K'$$

Type schemas
$$(\{K \ v\} \ \cup \ ctx) \ \Rightarrow \ \tau \implies (ctx \Rightarrow \tau)' \ [(v' :: s)/(v' :: (K' \cup s))]$$
$$\{\} \ \Rightarrow \ \tau \qquad\qquad \implies \tau'$$

Haskell type variables are translated to variables of class *pcpo*. Each type is associated to a sort in Isabelle (in Haskell, the same concept is called "kind"), defined by the set of the classes of which it is member.

Declarations
$$class \ K \ where \ (Dec_1; \ldots; Dec_n]) \ \implies \ class \ K' \ \subseteq \ pcpo; \ Dec'_1; \ \ldots; \ Dec'_n$$
$$f :: \phi \qquad\qquad\qquad \implies consts \ f' \ :: \ \phi'$$
$$type \ \tau \ = \ \tau_1 \qquad\qquad \implies \ type \ \tau \ = \ \tau'_1$$
$$(data \ \phi_1 \ = \ C_{11} \ x_1 \ldots x_i \ | \ \ldots \ | \ C_{1p} \ y_1 \ldots y_j;$$
$$\ldots;$$

Definitions
$$instance \ ctx \ \Rightarrow \ K_T \ (T \ v_1 \ \ldots \ v_n) \ where$$
$$(f_1 :: \tau_1 = t_1; \ldots; \ f_n :: \tau_n = t_n) \ \implies$$
$$instance$$
$$\tau' :: \ K'_T \ (\{pcpo\} \ \cup \ \{K' : (K \ v_1) \in ctx\}, \ \ldots,$$
$$\{pcpo\} \ \cup \ \{K' : (K \ v_n) \in ctx\})$$
$$with \ proof \ obligation;$$
$$defs \ f'_1 :: (ctx \Rightarrow \tau_1)' == t'_1; \ \ldots; \ f'_n :: (ctx \Rightarrow \tau_n)' == t'_n$$

Classes in Isabelle and Haskell are built quite differently. In Haskell, a type class is associated to a set of function declarations, and it can be interpreted

as the set of types where those functions are defined. In Isabelle, a type class has a single type parameter, it is associated to a set of axioms in a single type variable, and can be interpreted as the set of types that satisfy those axioms.

Not all the problems have been solved with respect to arities that may conflict in Isabelle, although they correspond to compatible Haskell instantiations. Moreover, Isabelle does neither allow for multi-parameter classes, nor for type constructor ones, therefore the same translation method cannot be applied to them.

Defined single-parameter classes are translated to Isabelle/HOLCF as subclasses of *pcpo* with empty axiomatization. Methods declarations associated with Haskell classes are translated to independent function declarations with appropriate class annotation on type variables. In Isabelle, each instance requires proofs that class axioms are satisfied by the instantiating type — anyway, as long as there are no axioms proofs are trivial and proof obligation may be automatically discharged. Method definitions associated with instance declarations are translated to independent function definitions, using type annotation and relying on Isabelle overloading.

In the internal representation of Haskell given by Programatica, function overloading is handled by means of dictionary parameters [Jon93]. This means that each function has additional parameters for the classes associated to its type variables. In fact, dictionary parameters are used to decide, for each instantiation of the function type variables, how to instantiate the methods called in the function body. On the other hand, overloading in Isabelle is obtained by adding explicitly type annotation to function definitions — dictionary parameters may thus be eliminated.

The translation of built-in classes may involve axioms — this is the case for equality. An Isabelle/HOLCF formalisation, based on the methods specification in [PJ03], has been given as follows in *HsHOLCF* (*neg* is lifted negation).

$consts$

$$heq :: \; 'a \; \to \; 'a \; \to \; tr$$
$$hneq :: 'a \; \to \; 'a \; \to \; tr$$

$axclass \; Eq \; < \; pcpo$
$eqAx : heq \cdot p \cdot q \; = \; neg \cdot (hneq \cdot p \cdot q)$

Functions *heq* and *hneq* can be defined, for each instantiating type, with the translation of equality and inequality, respectively. For each instance, a proof that the definitions satisfy *eqAx* needs to be given — the translation will simply print out *sorry* (a form of ellipsis in Isabelle). The definition of default methods for built-in types and the associated proofs can be found in *HsHOLCF*.

6 Translation of Function Definitions and Terms

Terms of built-in type are translated using Isabelle/HOLCF-defined lifting function *Def*. The bottom element $\perp$ is used for undefined terms.

Isabelle/HOLCF-defined $flift1 :: ('a \Rightarrow' b :: pcpo) \Rightarrow ('a \ lift \rightarrow' b)$ and $flift2 :: ('a \Rightarrow' b) \Rightarrow ('a \ lift \rightarrow' b \ lift)$ are used to lift operators, as well as the following, defined in *HsHOLCF*.

$$fliftbin :: ('a \Rightarrow \ 'b \Rightarrow \ 'c) \Rightarrow ('a \ lift \rightarrow \ 'b \ lift \rightarrow \ 'c \ lift)$$
$$fliftbin \ f \ == \ flift1 \ (\lambda x. \ flift2 \ (f \ x))$$

Boolean values are translated to values of *bool lift* (*tr* in Isabelle/HOLCF) i.e. *TT*, *FF* and $\bot$, and Boolean connectives to the corresponding Isabelle/HOLCF operators. Isabelle/HOLCF-defined *If then else fi* and *case* syntax are used to translate conditional and case expressions, respectively. There are restrictions, however, on case epressions, due to limitations in the translation of patterns; in particular, the case term has to be a variable, and only simple patterns are allowed (no nested ones). On the other hand, Isabelle sensitiveness to the order of patterns in case expressions is dealt with. Multiple function definitions are translated as definitions based on case expressions. In function definitions as well as in case expressions, both wildcards — not available in Isabelle — and incomplete patterns — not allowed — are dealt with by elimination, $\bot$ being used as default value in the latters. Only let expressions without patterns on the left are dealt with; where expressions, guarded expressions and list comprehension are not covered.

Definitions

$$f \ \overline{x} \ p_1 \ \overline{x}_1 \ = \ t_1; \ \ldots; \ f \ \overline{x} \ p_n \ \overline{x}_n \ = \ t_n \implies$$
$$\quad (f \ \overline{x} \ = \ case \ y \ of \ (p_1 \rightarrow (\backslash \overline{x}_1 \rightarrow \ t_1); \ \ldots; \ p_n(\rightarrow \backslash \overline{x}_n \rightarrow \ t_n)))'$$
$$f \ \overline{x} \ = \ t \qquad \implies defs \ f' :: \phi' \ == \ LAM \ \overline{x}'. \ t'$$
$$\quad with \ f :: \phi \ not \ occurring \ in \ t$$
$$(f_1 \ \overline{v_1} \ = \ t_1; \ \ldots; \ f_n \ \overline{v_n} \ = \ t_n) \implies$$
$$\quad fixrec \ f'_1 :: \phi'_1 \ = \ (LAM \ \overline{v_1}'. \ t'_1) \ and$$
$$\quad \ldots$$
$$\quad \quad and \ f'_n :: \phi'_n \ = \ (LAM \ \overline{v_n}'. \ t'_n)$$
$$\quad with \ f_1 :: \phi_1, \ldots, f_n :: \phi_n \ mutually \ recursive$$

Function declarations use Isabelle keyword *consts*. Datatype declarations in Isabelle/HOLCF are domain declarations and require explicitly destructors. Mutually recursive datatypes relies on specific Isabelle syntax (keyword *and*). Order of declarations is taken care of.

Non-recursive definitions are translated to standard definitions using Isabelle keyword *defs*. Recursive definitions rely on Isabelle/HOLCF package *fixrec* which provides nice syntax for fixed point definitions, including mutual recursion. Lambda abstraction is translated as continuous abstraction (LAM), function application as continuous application (the *dot* operator), equivalent to lambda abstraction (λ) and standard function application, respectively, when all arguments are continuous.

7 Example Proofs

8 Conclusion and future work

Concerning related work, although there have been translations of functional languages to first-order systems — those to FOL of Miranda [Tho94,Tho89,HT95] and Haskell [Tho92], both based on large-step operational semantics; that of Haskell to Agda implementation of Martin-Loef type theory in [ABB$^+$05] — still, higher-order logic may be quite helpful in order to deal with features such as currying and polymorphism. Moreover, higher-order approaches may rely on denotational semantics — as for examples, [HMW05] translating Haskell to HOLCF, and [LP04] translating ML to HOL — allowing for program representation closer to specification as well as for proofs comparatively more abstract and general.

The translation of Haskell to Isabelle/HOLCF proposed in [HMW05] uses deep embedding to deal with types. Haskell types are translated to terms, relying on a domain-theoretic modelling of the type system at the object level, allowing explicitly for a clear semantics, and providing for an implementation that can capture most features, including type constructor classes. In contrast, we provide in the case of Isabelle/HOLCF with a translation that follows the lines of a denotational semantics under the assumption that type constructors and type application in Haskell can be mapped to corresponding constructors and built-in application in Isabelle without loss from the point of view of behavioural equivalence between programs — in particular, translating Haskell datatypes to Isabelle ones. Our solution gives in general less expressiveness than the deeper approach — however, when we can get it to deal with cases of interest, it might make proofs easier.

Isabelle does not allow for type constructor classes, therefore there is hardly a way shallow embedding of Haskell types may extend to cover them. This limitation is particularly acute with respect to monads and *do* notation. The problem is brilliantly avoided in [HMW05] by resorting to a deeper modelling of types. operator.

The main advantage of shallow embedding is to get as much as possible out of the automation currently available in Isabelle, especially with respect to type checking. Isabelle/HOLCF in particular provides with an expressive semantics covering lazy evaluation, as well as with a smart syntax — also thanks to the *fixrec* package. The main disadvantage lies with lack of type constructor classes. Anyway, it is possible to get around the obstacle, at least partially, by relying on an axiomatic characterisation of monads and on a proof-reuse strategy that actually minimises the need for interactive proofs.

Future work should use this framework for proving properties of Haskell programs. For monadic programs, we are also planning to use the monad-based dynamic Hoare and dynamic logic that already have been formalised in Isabelle [Wal05]. Our translation tool from Haskell to Isabelle is part of the Heterogeneous Tool Set Hets and can be downloaded from `http://www.dfki.de/sks/hets`. More details about the translations can be found in [TLMM07].

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